



# Advance Database -Lecture 3

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SEVENTH EDITION

# Database System Concepts



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# Chapter 18 : Concurrency Control

**Database System Concepts, 7<sup>th</sup> Ed.**

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# Outline

- Lock-Based Protocols
- Timestamp-Based Protocols
- Validation-Based Protocols
- Multiple Granularity
- Multiversion Schemes
- Insert and Delete Operations
- Concurrency in Index Structures



# Lock-Based Protocols

- A lock is a mechanism to control concurrent access to a data item
- Data items can be locked in two modes :
  1. **exclusive** (*X*) *mode*. Data item can be both read as well as written. X-lock is requested using **lock-X** instruction.
  2. **shared** (*S*) *mode*. Data item can only be read. S-lock is requested using **lock-S** instruction.
- Lock requests are made to concurrency-control manager. Transaction can proceed only after request is granted.



# Lock-Based Protocols (Cont.)

- **Lock-compatibility matrix**

	S	X
S	true	false
X	false	false

- A transaction may be granted a lock on an item if the requested lock is compatible with locks already held on the item by other transactions
- Any number of transactions can hold shared locks on an item,
- But if any transaction holds an exclusive on the item no other transaction may hold any lock on the item.



## Lock-Based Protocols (Cont.)

- Example of a transaction performing locking:

$T_2$ : **lock-S**(A);

**read** (A);

**unlock**(A);

**lock-S**(B);

**read** (B);

**unlock**(B);

**display**(A+B)

- Locking as above is not sufficient to guarantee serializability



- Suppose that the values of accounts A and B are \$100 and \$200,. If these two transactions are executed serially, order T1, T2 or the order T2, T1,
- Then transaction T2 will display the value \$300.

$T_1$ : lock-X( $B$ );  
read( $B$ );  
 $B := B - 50$ ;  
write( $B$ );  
unlock( $B$ );  
lock-X( $A$ );  
read( $A$ );  
 $A := A + 50$ ;  
write( $A$ );  
unlock( $A$ ).

$T_2$ : lock-S( $A$ );  
read( $A$ );  
unlock( $A$ );  
lock-S( $B$ );  
read( $B$ );  
unlock( $B$ );  
display( $A + B$ ).



# Schedule With Lock Grants

- Grants omitted in rest of chapter
  - Assume grant happens just before the next instruction following lock request
- This schedule is not serializable (why?)
- A **locking protocol** is a set of rules followed by all transactions while requesting and releasing locks.
- Locking protocols enforce serializability by restricting the set of possible schedules.

$T_1$	$T_2$	concurrency-control manager
lock-X( $B$ )		grant-X( $B, T_1$ )
read( $B$ ) $B := B - 50$ write( $B$ ) unlock( $B$ )	lock-S( $A$ )	grant-S( $A, T_2$ )
	read( $A$ ) unlock( $A$ ) lock-S( $B$ )	grant-S( $B, T_2$ )
	read( $B$ ) unlock( $B$ ) display( $A + B$ )	
lock-X( $A$ )		grant-X( $A, T_1$ )
read( $A$ ) $A := A + 50$ write( $A$ ) unlock( $A$ )		





# Schedule With Lock Grants

- Same schedule but delayed unlocks:

$T_1$	$T_2$
lock-X( $B$ );	lock-S( $A$ );
read( $B$ );	read( $A$ );
$B := B - 50$ ;	lock-S( $B$ );
write( $B$ );	read( $B$ );
lock-X( $A$ );	display( $A + B$ );
read( $A$ );	unlock( $A$ );
$A := A + 50$ ;	unlock( $B$ ).
write( $A$ );	
unlock( $B$ );	
unlock( $A$ ).	



# Deadlock

$T_3$	$T_4$
lock-X( $B$ )	
read( $B$ )	
$B := B - 50$	
write( $B$ )	
	lock-S( $A$ )
	read( $A$ )
	lock-S( $B$ )
lock-X( $A$ )	

- Consider the partial schedule
- Neither  $T_3$  nor  $T_4$  can make progress — executing **lock-S( $B$ )** causes  $T_4$  to wait for  $T_3$  to release its lock on  $B$ , while executing **lock-X( $A$ )** causes  $T_3$  to wait for  $T_4$  to release its lock on  $A$ .
- Such a situation is called a **deadlock**.
  - To handle a deadlock one of  $T_3$  or  $T_4$  must be rolled back and its locks released.



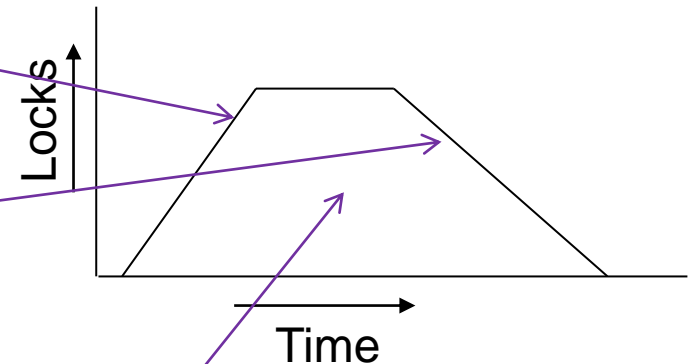
## Deadlock (Cont.)

- The potential for deadlock exists in most locking protocols. Deadlocks are a necessary evil.
- **Starvation** is also possible if concurrency control manager is badly designed. For example:
  - A transaction may be waiting for an X-lock on an item, while a sequence of other transactions request and are granted an S-lock on the same item.
  - The same transaction is repeatedly rolled back due to deadlocks.
- Concurrency control manager can be designed to prevent starvation.



# The Two-Phase Locking Protocol

- A protocol which ensures conflict-serializable schedules.
- Phase 1: **Growing Phase**
  - Transaction may **obtain locks**
  - Transaction may **not release locks**
- Phase 2: **Shrinking Phase**
  - Transaction may **release locks**
  - Transaction may **not obtain locks**
- The protocol **assures serializability**. It can be proved that the transactions can be **serialized** in the order of their **lock points** (i.e., the point where a transaction acquired its final lock).





# Cascading rollback may occur under two-phase locking.

- Each transaction observes the two-phase locking protocol, but the failure of T5 after the **read(A)** step of T7 leads to cascading rollback of T6 and T7

$T_5$	$T_6$	$T_7$
lock-X(A) read(A) lock-S(B) read(B) write(A) unlock(A)	lock-X(A) read(A) write(A) unlock(A)	lock-S(A) read(A) ←



# The Two-Phase Locking Protocol (Cont.)

- Two-phase locking *does not ensure* freedom from **deadlocks**
- Extensions to basic **two-phase locking** needed to ensure **recoverability** of freedom from **cascading roll-back**
  - **Strict two-phase locking**: a transaction must **hold all** its **exclusive locks** till it **commits/aborts**.
    - Ensures **recoverability** and **avoids** cascading **roll-backs**
  - **Rigorous two-phase locking**: a transaction must **hold all locks** till **commit/abort**.
    - Transactions can be serialized in the order in which they commit.
- Most **databases** implement **rigorous** two-phase locking, *but refer to it as simply two-phase locking*



# The Two-Phase Locking Protocol (Cont.)

- Two-phase locking is **not** a **necessary** condition for **serializability**
  - There are **conflict serializable schedules** that cannot be obtained if the **two-phase locking protocol** is used.
- In the absence of extra information (e.g., **ordering of access to data**), two-phase locking is necessary for **conflict serializability** *in the following sense*:
  - *Given a transaction  $T_i$  that does not follow two-phase locking, we can find a transaction  $T_j$  that uses two-phase locking, and a schedule for  $T_i$  and  $T_j$  that is not conflict serializable.*



# Locking Protocols

- Given a locking protocol (such as 2PL)
  - A schedule  $S$  is **legal** under a locking protocol if it can be generated by a set of **transactions** that **follow the protocol**
  - A **protocol ensures** serializability if **all legal schedules** under that **protocol are serializable**





# Lock Conversions

- Two-phase locking protocol with **lock conversions**:
  - **Growing Phase:**
    - can acquire a lock-S on item
    - can acquire a lock-X on item
    - can **convert** a lock-S to a lock-X (**upgrade**)
  - **Shrinking Phase:**
    - can release a lock-S
    - can release a lock-X
    - can convert a lock-X to a lock-S (**downgrade**)
- This protocol ensures serializability



# Automatic Acquisition of Locks

- A transaction  $T_i$  issues the standard read/write instruction, without explicit locking calls.
- The operation **read( $D$ )** is **processed** as:

**if**  $T_i$  has a **lock** on  $D$

**then**

    read( $D$ )

**else begin**

    if necessary wait until no other  
    transaction has a **lock-X** on  $D$

    grant  $T_i$  a **lock-S** on  $D$ ;

    read( $D$ )

**end**



# Automatic Acquisition of Locks (Cont.)

- The operation **write**( $D$ ) is processed as:  
if  $T_i$  has a **lock-X** on  $D$   
  **then**  
    **write**( $D$ )  
  **else begin**  
    if necessary wait until no other trans. has any lock on  $D$ ,  
    if  $T_i$  has a **lock-S** on  $D$   
      **then**  
        **upgrade** lock on  $D$  to **lock-X**  
      **else**  
        grant  $T_i$  a **lock-X** on  $D$   
    **write**( $D$ )  
  **end;**
- **All locks are released after commit or abort**



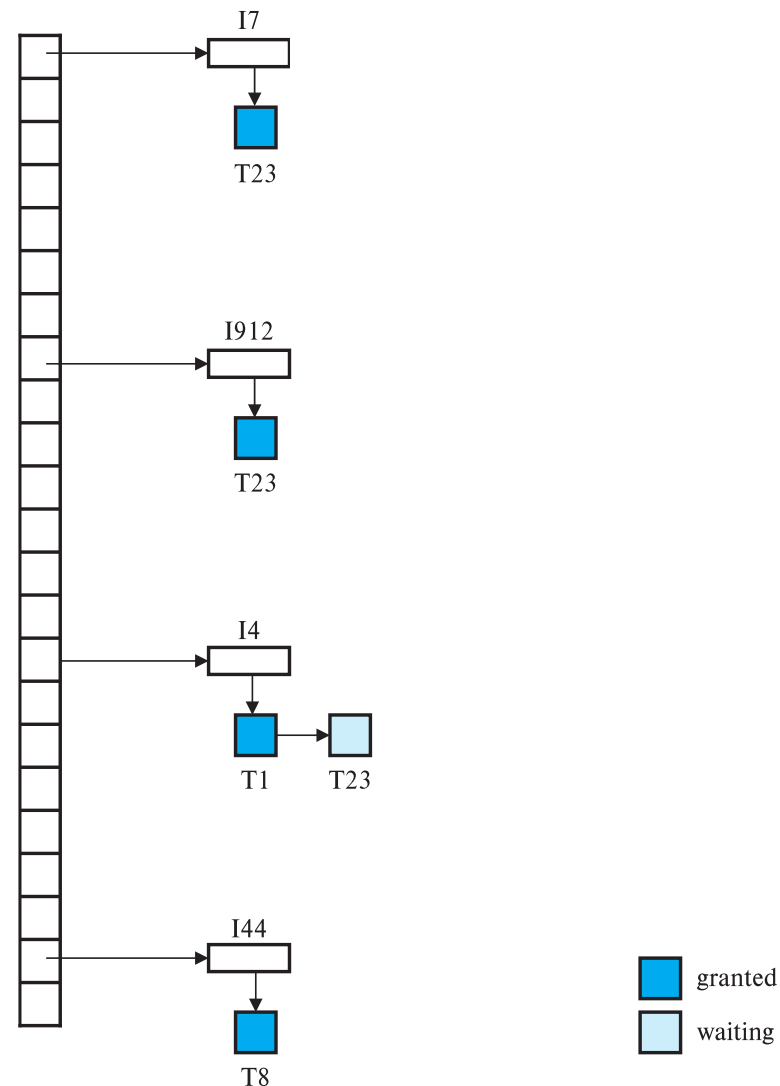
# Implementation of Locking

- A **lock manager** can be implemented as a separate process
- Transactions can **send lock** and **unlock** requests as **messages**
- The **lock manager replies** to a **lock request** by **sending a lock grant** messages (or a message **asking the transaction to roll back, in case of a deadlock**)
  - The requesting transaction waits until its request is answered
- The **lock manager** maintains an in-memory data-structure called a **lock table** to record **granted** locks and **pending** requests



# Lock Table

- **Dark rectangles** indicate **granted locks**, **light colored** ones indicate **waiting requests**
- Lock table also records the **type of lock** granted or requested
- New request is added to the end of the queue of requests for the data item, and granted if it is compatible with all earlier locks
- **Unlock requests** result in the **request** being **deleted**, and later **requests** are checked to see if they can now be **granted**
- If **transaction aborts**, all **waiting** or **granted** requests of the transaction are **deleted**
  - **lock manager** may keep a **list of locks** held by **each transaction**, to implement this efficiently





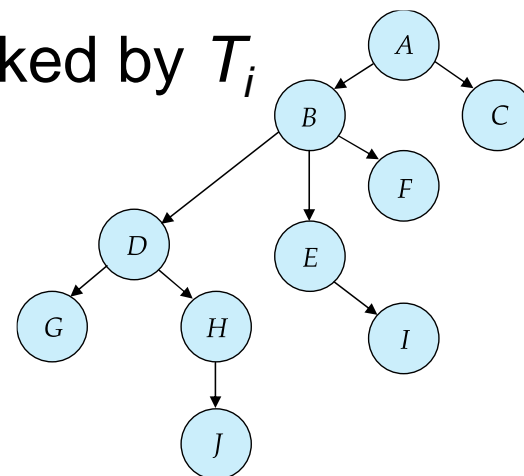
# Graph-Based Protocols

- Graph-based protocols are an alternative to two-phase locking
- Impose a partial ordering  $\rightarrow$  on the set  $\mathbf{D} = \{d_1, d_2, \dots, d_h\}$  of all data items.
  - If  $d_i \rightarrow d_j$  then any transaction accessing both  $d_i$  and  $d_j$  must access  $d_i$  before accessing  $d_j$ .
  - Implies that the set  $\mathbf{D}$  may now be viewed as a directed acyclic graph, called a *database graph*.
- The *tree-protocol* is a simple kind of graph protocol.



# Tree Protocol

- In the tree protocol, the only lock instruction allowed is **exclusive locks (lock-X)**.
- Each transaction  $T_i$  can lock a data item at most **once**, and must observe the following rules:
  1. The **first lock** by  $T_i$  may be on **any data item**. **Subsequently, a data  $Q$  can be locked by  $T_i$  only if the parent of  $Q$  is currently locked by  $T_i$ .**
  2. Data items may be unlocked at any time.
  3. A data item that has been locked and unlocked by  $T_i$  cannot subsequently be relocked by  $T_i$ .





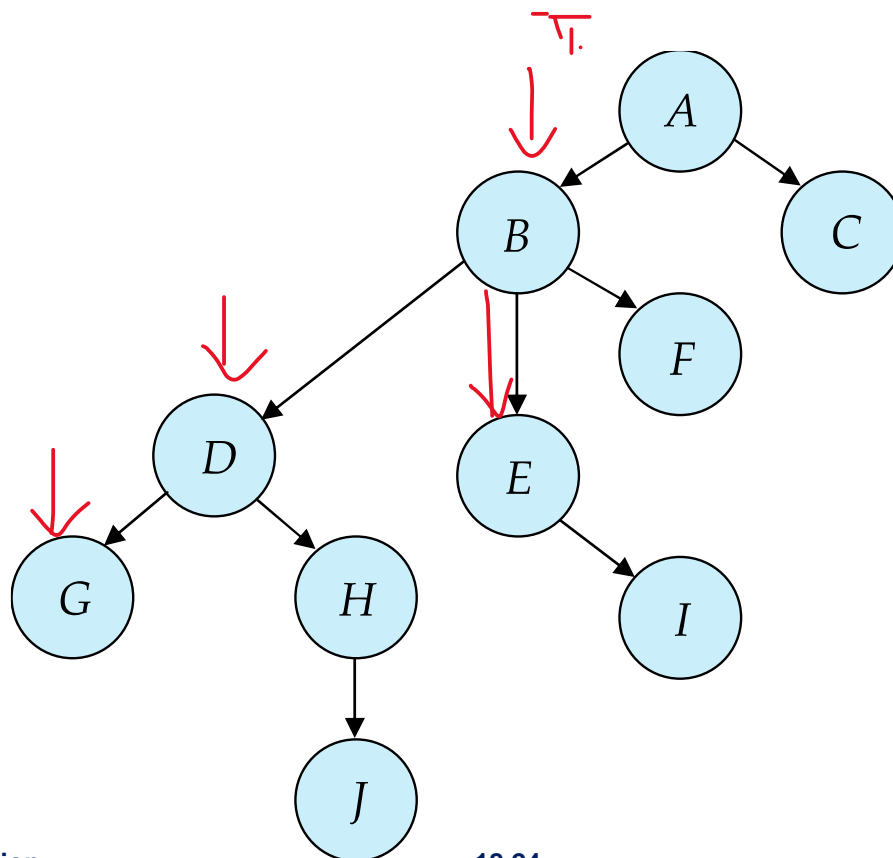
# Tree Protocol

$T_{10}$ : lock-X( $B$ ); lock-X( $E$ ); lock-X( $D$ ); unlock( $B$ ); unlock( $E$ ); lock-X( $G$ );  
unlock( $D$ ); unlock( $G$ ).

$T_{11}$ : lock-X( $D$ ); lock-X( $H$ ); unlock( $D$ ); unlock( $H$ ).

$T_{12}$ : lock-X( $B$ ); lock-X( $E$ ); unlock( $E$ ); unlock( $B$ ).

$T_{13}$ : lock-X( $D$ ); lock-X( $H$ ); unlock( $D$ ); unlock( $H$ ).







# Tree Protocol

- One possible schedule in which these four transactions participated appears in Figure
- During its execution, transaction T10 holds locks on two disjoint subtrees. Observe that the schedule of Figure is conflict serializable.
- It can be shown not only that the tree protocol ensures conflict serializability, but also that this protocol ensures freedom from deadlock

$T_{10}$	$T_{11}$	$T_{12}$	$T_{13}$
lock-X( $B$ )	lock-X( $D$ ) lock-X( $H$ ) unlock( $D$ )		
lock-X( $E$ ) lock-X( $D$ ) unlock( $B$ ) unlock( $E$ )		lock-X( $B$ ) lock-X( $E$ )	
lock-X( $G$ ) unlock( $D$ )	unlock( $H$ )		
			lock-X( $D$ ) lock-X( $H$ ) unlock( $D$ ) unlock( $H$ )
unlock( $G$ )		unlock( $E$ ) unlock( $B$ )	



## Graph-Based Protocols (Cont.)

- The tree protocol **ensures conflict serializability** as well as **freedom from deadlock**.
- **Unlocking** may occur **earlier** in the **tree-locking** protocol than in the **two-phase locking** protocol.
  - **Shorter waiting times**, and **increase in concurrency**
  - Protocol is **deadlock-free**, **no rollbacks** are required



## Graph-Based Protocols (Cont.)

### ■ Drawbacks

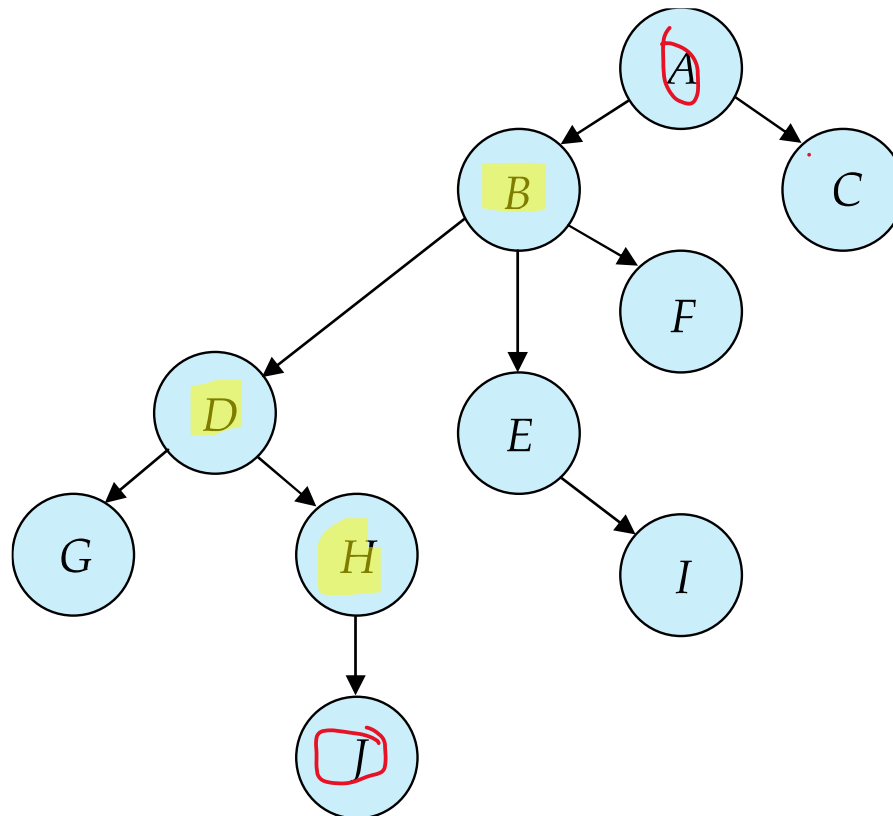
- Protocol does **not guarantee** recoverability or **cascade** freedom
  - Need to introduce **commit dependencies** to **ensure recoverability**
- **Transactions** may have to **lock data items** that they **do not access**.
  - increased **locking overhead**, and additional **waiting** time
  - potential **decrease** in **concurrency**
- Schedules **not possible** under **two-phase locking** are **possible** under the **tree protocol**, and **vice versa**.



# Graph-Based Protocols (Cont.)

## ■ Drawback:

- a transaction that needs to access data items **A** and **J** in the database graph of must lock not only A and J, but also data items **B, D, H**





# Deadlock Handling

- System is **deadlocked** if there is a set of transactions such that every transaction in the set is waiting for another transaction in the set.

$T_3$	$T_4$
lock-X( $B$ )	
read( $B$ )	
$B := B - 50$	
write( $B$ )	
	lock-S( $A$ )
	read( $A$ )
	lock-S( $B$ )
lock-X( $A$ )	



# Deadlock Handling

- ***Deadlock prevention*** protocols **ensure** that the system will ***never*** enter into a **deadlock** state. Some prevention strategies:
  - Require that **each transaction locks** all its **data items** before it **begins execution** (pre-declaration).
  - Impose **partial ordering** of all data items and require that a **transaction** can lock data items only in the **order specified** by the **partial order** (graph-based protocol).



# More Deadlock Prevention Strategies

- **wait-die** scheme — non-preemptive
  - **Older** transaction may **wait for younger** one to release data item.
  - **Younger** transactions **never wait** for **older ones**; they are rolled back instead.
  - A transaction may die several times before acquiring a lock



## wait-die scheme, Example

- suppose that transactions **T14**, **T15**, and **T16** have timestamps **5**, **10**, and **15**, respectively.
- If T14 requests a data item held by T15, then what will happen?
- If T16 requests a data item held by T15, then what will happen?





# More Deadlock Prevention Strategies

- **wound-wait** scheme — **preemptive**
  - **Older** transaction *wounds* (**forces rollback**) of **younger transaction** instead of waiting for it.
  - **Younger** transactions may **wait for older ones**.
  - **Fewer** rollbacks than *wait-die* scheme.
- In both schemes, a **rolled back transactions** is restarted with **its original timestamp**.
  - Ensures that **older** transactions have **precedence** over **newer ones**, and **starvation** is thus **avoided**.



## Wound-die scheme, Example

- suppose that transactions **T14**, **T15**, and **T16** have timestamps **5**, **10**, and **15**, respectively.
- if T14 requests a data item held by T15, then what happens?
- If T16 requests a data item held by T15, then what happens?



# Deadlock prevention (Cont.)

## ■ Timeout-Based Schemes:

- A transaction **waits** for a lock only for a **specified amount of time**. After that, the **wait times out** and the transaction is **rolled back**.
- Ensures that **deadlocks** get **resolved** by **timeout** if they occur
- Simple to implement
- But may roll back transaction unnecessarily in absence of deadlock
  - Difficult to **determine** good value of the **timeout interval**.
- **Starvation** is also **possible**



# Deadlock Detection

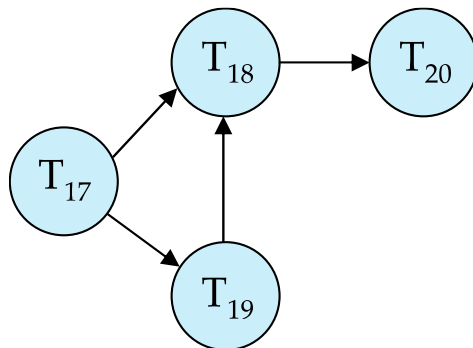
in the systems that do not prevent deadlock

- **Wait-for graph**

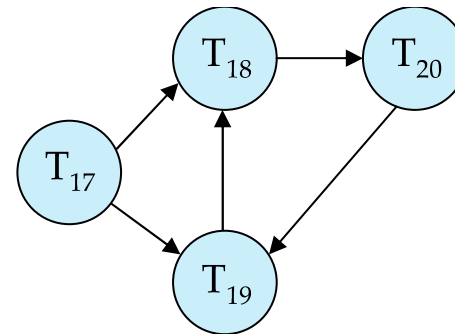
- *Vertices: transactions*
- *Edge from  $T_i \rightarrow T_j$  : if  $T_i$  is waiting for a lock held in conflicting mode by  $T_j$*

- The system is in a deadlock state if and only if the **wait-for graph has a cycle**.

- **Invoke a deadlock-detection** algorithm periodically to look for cycles.



Wait-for graph without a cycle



Wait-for graph with a cycle



# Deadlock Recovery

- When **deadlock** is **detected** :
  - Some **transaction** will have to **rolled back** (made a **victim**) to **break deadlock cycle**.
    - **Select** that **transaction** as victim that will incur **minimum cost**
  - Rollback -- determine **how far to roll back** transaction
    - **Total rollback**: **Abort** the transaction and then restart it.
    - **Partial rollback**: **Rollback** victim transaction only as far as necessary to release locks that another transaction in cycle is waiting for
- **Starvation** can **happen** (why?)
  - One solution: **oldest transaction** in the **deadlock** set is **never chosen as victim**



# Multiple Granularity

- it would be advantageous to group several data items, and to treat them as one individual synchronization unit.
- For example, if a **transaction  $T_i$**  needs to access an **entire relation**, and a locking protocol is used to **lock tuples**, then  $T_i$  must **lock each tuple in the relation**.
  - **acquiring many such locks is time-consuming**; even worse, the **lock table may become very large** and no longer fit in memory.
  - It would be better if  $T_i$  could issue a **single lock request** to lock the **entire relation**.
    - if transaction  $T_i$  needs to access only a few tuples, it should not be required to lock the entire relation,
- a mechanism to allow the system to define **multiple levels of granularity is needed then**



# Multiple Granularity

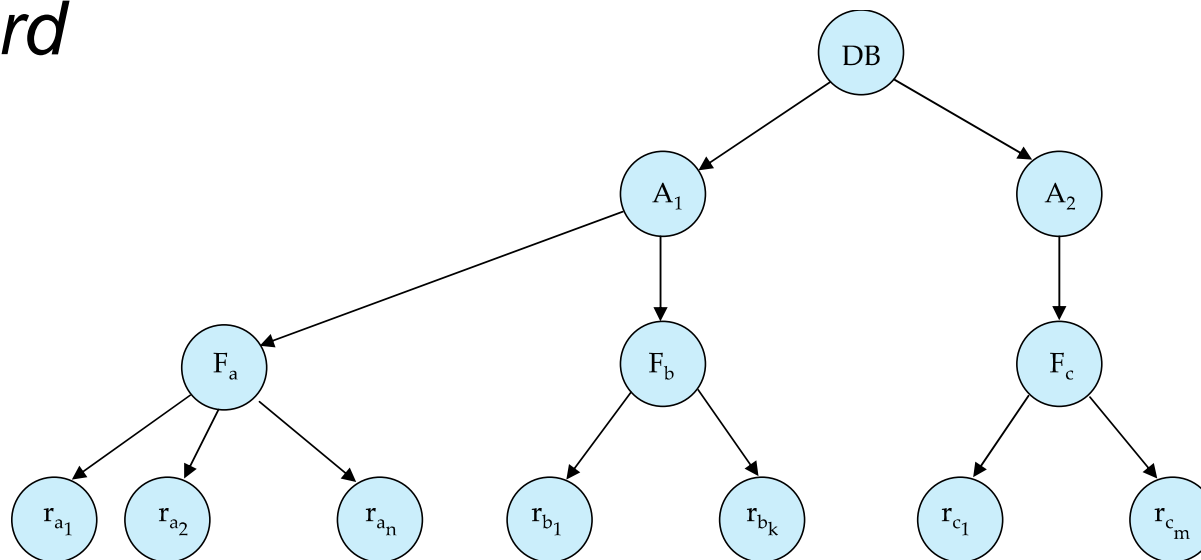
- Allow **data items to be of various sizes** and define a **hierarchy of data granularities**, where the small granularities are nested within **larger ones**
- Can be represented graphically as a tree (but do not confuse with tree-locking protocol)
- When a transaction **locks a node in the tree *explicitly***, it *implicitly* locks all the **node's descendants** in the same mode.
- **Granularity of locking** (level in tree where locking is done):
  - **Fine granularity** (lower in tree): high concurrency, high locking overhead
  - **Coarse granularity** (higher in tree): low locking overhead, low concurrency



# Example of Granularity Hierarchy

The levels, starting from the coarsest (top) level are

- *database*
- *area*
- *file*
- *record*

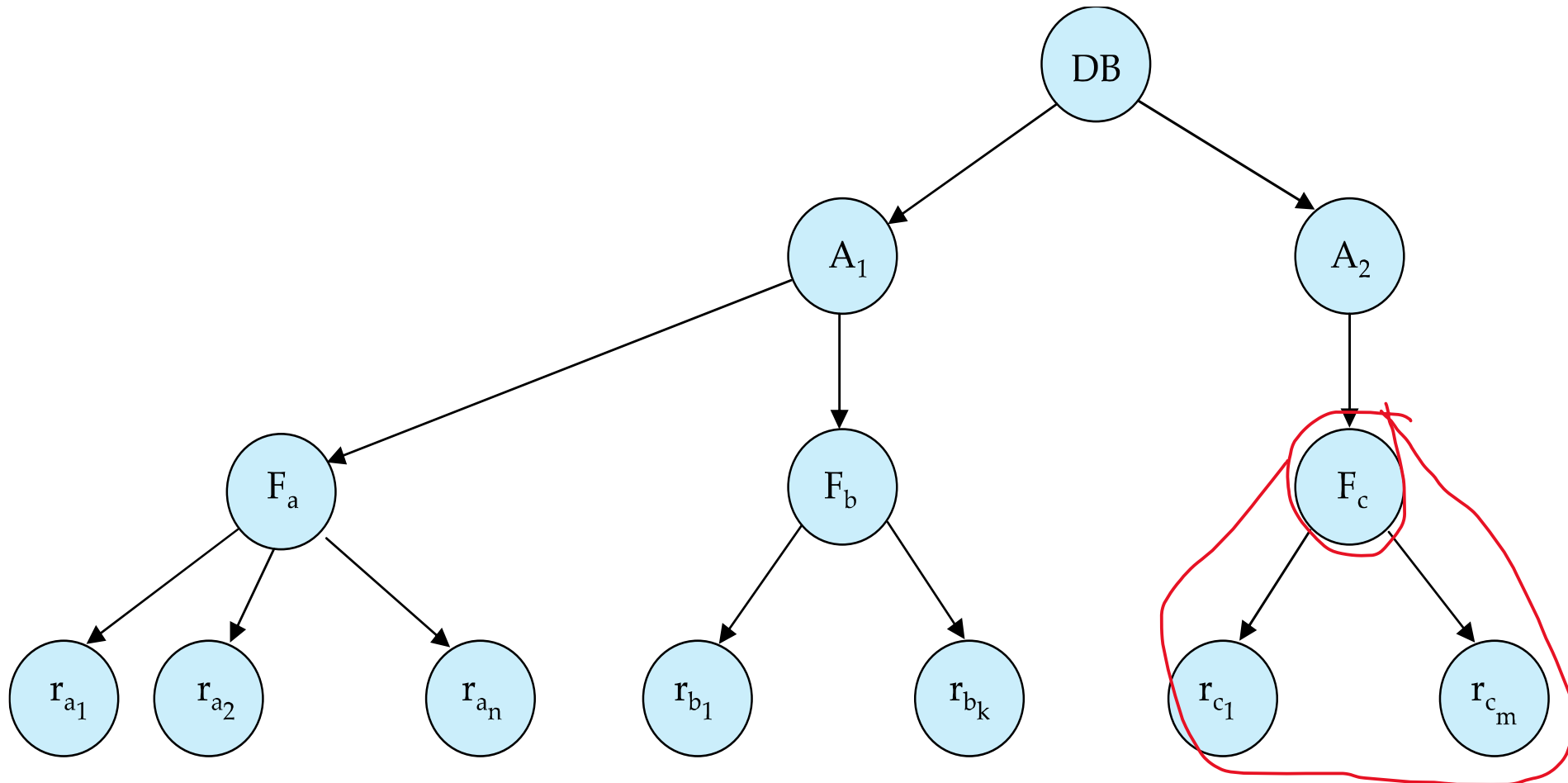






# Example of Granularity Hierarchy

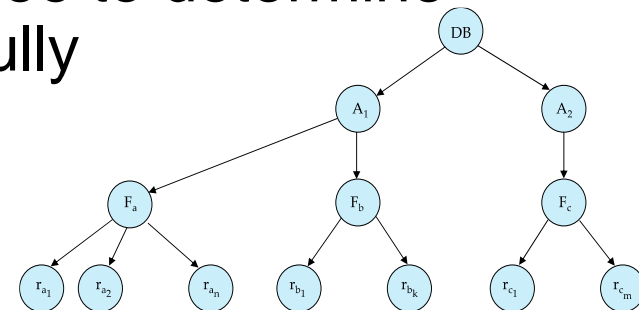
- if transaction  $T_i$  gets an explicit lock on file  $F_c$  of Figure, in exclusive mode, then it has an **implicit lock in exclusive mode** on all the records belonging to that file.





# Example of Granularity Hierarchy

- how does the system determine if the root node can be locked?
  - One possibility is for it to search the entire tree.
    - This solution, is not suitable
  - A more efficient way: a new class of lock modes, called intention lock modes.
    - If a node is locked in an intention mode, explicit locking is done at a lower level of the tree (that is, at a finer granularity)
    - **Intention** locks are put on all the **ancestors** of a node before that node is locked explicitly. So a transaction does not need to search the entire tree to determine whether it can lock a node successfully





# Intention Lock Modes

- In addition to **S** and **X** lock modes, there are three **additional lock** modes with multiple granularity:
  - ***intention-shared* (IS)**: indicates **explicit locking at a lower level of the tree** but only with **shared locks**.
  - ***intention-exclusive* (IX)**: indicates explicit locking at a lower level with **exclusive or shared** locks
  - ***shared and intention-exclusive* (SIX)**: if a node is locked in (SIX) mode, the **subtree rooted** by that node is **locked explicitly in shared mode** and **explicit** locking is being done at a lower level with **exclusive-mode locks**.
- **Intention** locks allow a **higher level node to be locked** in **S** or **X** mode without having to check all descendent nodes.



# Compatibility Matrix with Intention Lock Modes

- The compatibility matrix for all lock modes is:

	IS	IX	S	SIX	X
IS	true	true	true	true	false
IX	true	true	false	false	false
S	true	false	true	false	false
SIX	true	false	false	false	false
X	false	false	false	false	false



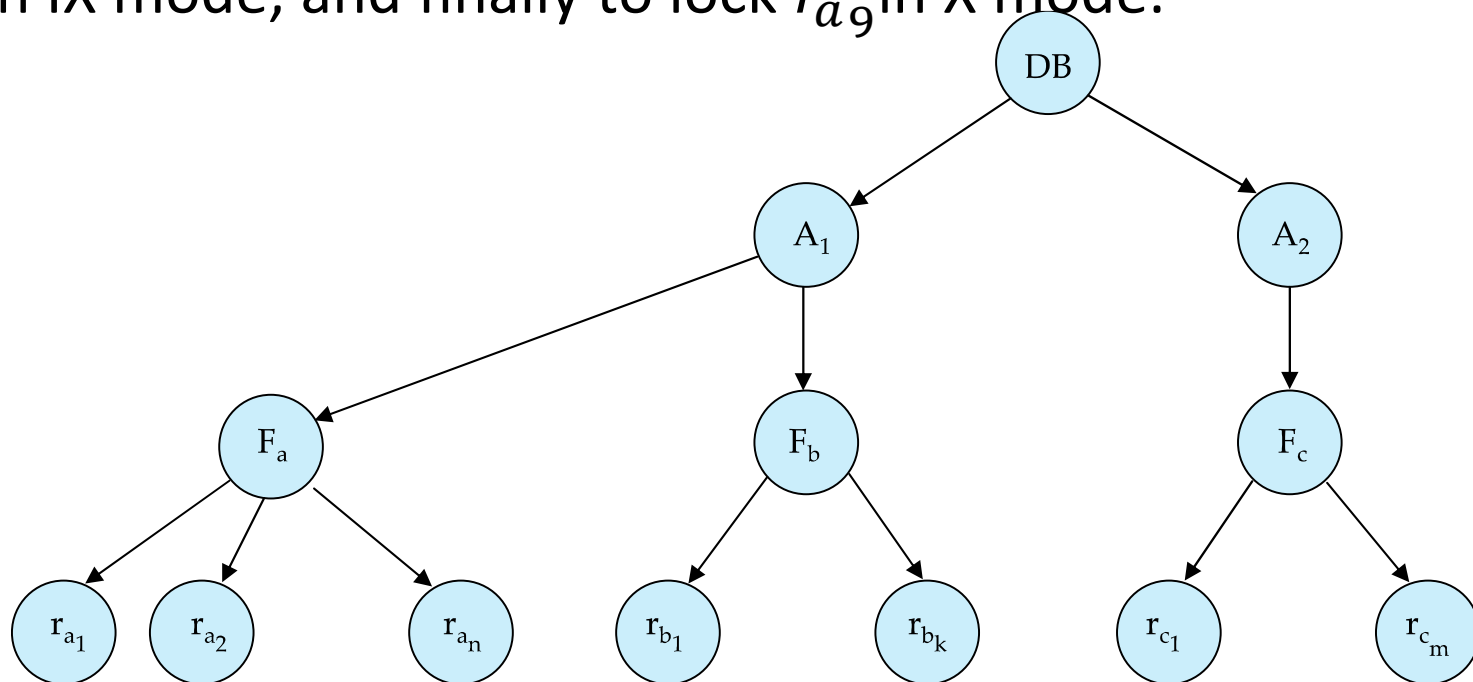
# Multiple Granularity Locking Scheme

- Transaction  $T_i$  can lock a node  $Q$ , using the following rules:
  1. The **lock compatibility** matrix must be observed.
  2. The root of the tree must be **locked first**, and may be locked in any mode.
  3. A node  $Q$  can be locked by  $T_i$  in S or IS mode only if the **parent** of  $Q$  is currently locked by  $T_i$  in either **IX or IS** mode.
  4. A node  $Q$  can be locked by  $T_i$  in **X, SIX, or IX** mode only if the parent of  $Q$  is currently locked by  $T_i$  in either **IX or SIX** mode.
  5.  $T_i$  can lock a node only if it has not previously unlocked any node (that is,  $T_i$  is two-phase).
  6.  $T_i$  can unlock a node  $Q$  only if none of the children of  $Q$  are currently locked by  $T_i$ .
- Observe that locks are acquired in root-to-leaf order, whereas they are released in leaf-to-root order.
- **Lock granularity escalation**: in case there are too many locks at a particular level, switch to higher granularity S or X lock



# Multiple Granularity Locking Scheme

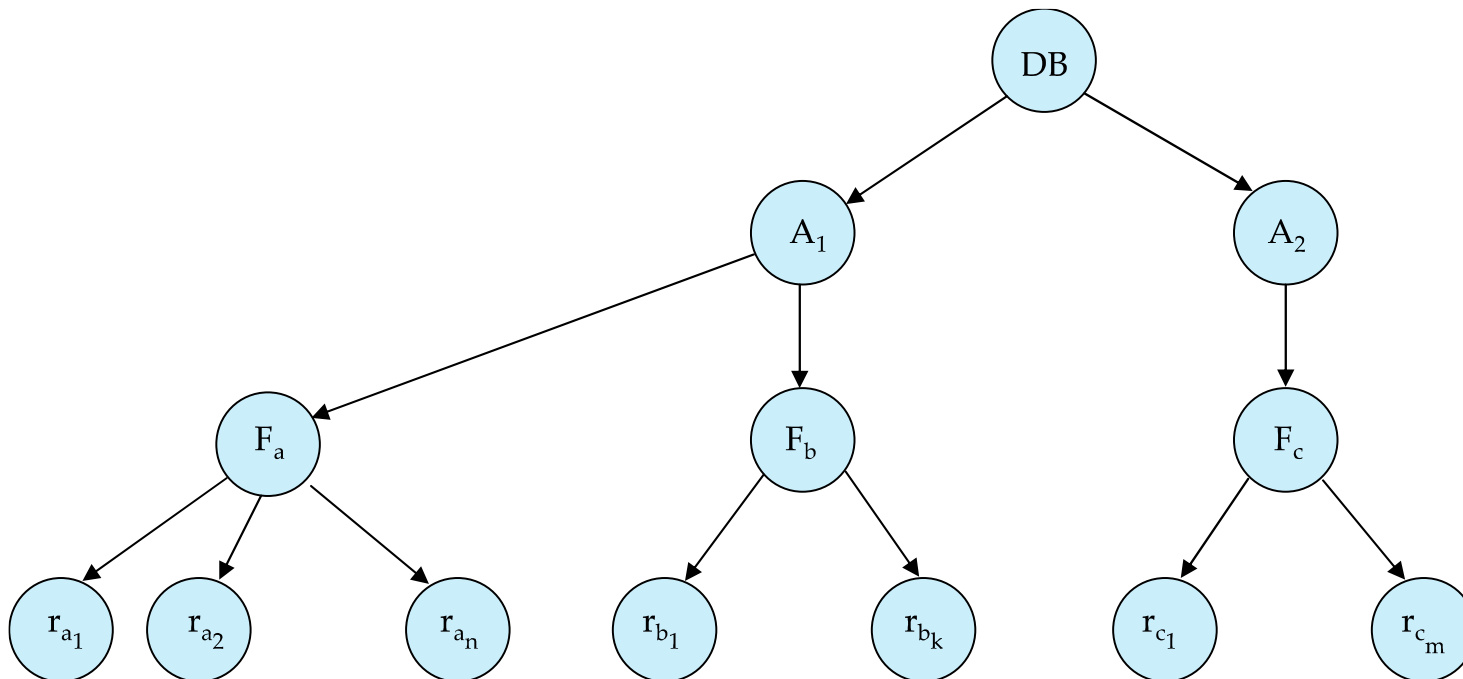
- Suppose that transaction **T21** reads record  $r_{a_2}$  in file  $F_a$ . Then, T21 needs to lock the database, area  $A_1$ , and  $F_a$  in IS mode (and in that order), and finally to lock  $r_{a_2}$  in S mode.
- Suppose that transaction T22 modifies record  $r_{a_9}$  in file  $F_a$ . Then, T22 needs to lock the database, area  $A_1$ , and file  $F_a$  (and in that order) in IX mode, and finally to lock  $r_{a_9}$  in X mode.





# Multiple Granularity Locking Scheme

- Suppose that transaction **T23** reads all the records in file  $F_a$ . Then, T23 needs to lock the database and area A1 (and in that order) in IS mode, and finally to lock  $F_a$  in S mode.
- Suppose that transaction T24 reads the entire database. It can do so after locking the database in S mode.





# Timestamp Based Concurrency Control





# Timestamp Based Concurrency Control

- The locking protocols that we have described thus far determine the order between every pair of conflicting transactions at execution time by the first lock that both members of the pair request that involves incompatible modes.
- a timestamp-ordering scheme
  - For **determining the serializability order** is to select an **ordering among transactions in advance**.
  - With each transaction  $T_i$  in the system, **we associate a unique fixed timestamp**, denoted by  **$TS(T_i)$** . This timestamp is assigned by the database system before the transaction  $T_i$  starts execution



# Timestamp-Based Protocols

- Each **transaction**  $T_i$  is issued a **timestamp**  $TS(T_i)$  when it enters the system.
  - **Each** transaction has a **unique** timestamp
  - **Newer** transactions have timestamps strictly **greater** than earlier ones  $TS(T_i) < TS(T_j)$ 
    - Timestamp could be based on a **logical counter**
    - system clock as the timestamp;
  - Timestamp-based protocols manage concurrent execution such that  
**time-stamp order = serializability order.**
    - **The system must ensure it is equivalent to a serial Schedule**



# Timestamp-Ordering Protocol

## The **timestamp ordering (TSO)** protocol

- Maintains for each data **Q** **two timestamp** values:
  - **W-timestamp(Q)** is the largest time-stamp of any transaction that executed **write(Q)** successfully.
  - **R-timestamp(Q)** is the **largest time-stamp** of any transaction that executed **read(Q)** successfully.
- Imposes rules on read and write operations to ensure that
  - Any **conflicting operations** are executed in **timestamp order**
  - **Out of order** operations cause transaction rollback



## Timestamp-Based Protocols (Cont.)

- Suppose a transaction  $T_i$  issues a **read**( $Q$ )
  1. If  $\mathbf{TS}(T_i) < \mathbf{W}$ -timestamp( $Q$ ), then  $T_i$  needs to read a value of  $Q$  that was already **overwritten**.
    - Hence, the **read** operation is rejected, and  $T_i$  is rolled back.
  2. If  $\mathbf{TS}(T_i) \geq \mathbf{W}$ -timestamp( $Q$ ), then the **read** operation is executed, and R-timestamp( $Q$ ) is set to  $=\mathbf{max}(\mathbf{R}$ -timestamp( $Q$ ),  $\mathbf{TS}(T_i))$ .



# Timestamp-Based Protocols (Cont.)

- Suppose that transaction  $T_i$  issues **write**( $Q$ ).
  1. If  $TS(T_i) < R\text{-timestamp}(Q)$ , then the value of  $Q$  that  $T_i$  is producing was needed previously, and the system assumed that that value would never be produced.
    - Hence, the **write** operation is rejected, and  $T_i$  is rolled back.
  2. If  $TS(T_i) < W\text{-timestamp}(Q)$ , then  $T_i$  is attempting to write an **obsolete** value of  $Q$ .
    - Hence, this **write** operation is rejected, and  $T_i$  is rolled back.
  3. Otherwise, the **write** operation is executed, and  $W\text{-timestamp}(Q)$  is set to  $TS(T_i)$ .



## Timestamp-Based Protocols (Cont.)

- If a transaction  $T_i$  is rolled back by the concurrency-control scheme as result of issuance of either a read or write operation, the system assigns it a new timestamp and restarts it.



# Example of Schedule Under TSO

- Is this schedule valid under TSO?
  - We consider transactions  $T_{25}$  and  $T_{26}$ . Transaction  $T_{25}$  displays the contents of accounts  $A + B$
  - Transaction  $T_{26}$  transfers \$50 from account  $B$  to account  $A$ , and then displays the contents of both

$T_{25}$	$T_{26}$
read( $B$ )	read( $B$ ) $B := B - 50$ write( $B$ )
read( $A$ )	read( $A$ )
display( $A + B$ )	$A := A + 50$ write( $A$ ) display( $A + B$ )

Assume that initially:

$$R\text{-TS}(A) = W\text{-TS}(A) = 0$$

$$R\text{-TS}(B) = W\text{-TS}(B) = 0$$

Assume  $TS(T_{25}) = 25$  and

$$TS(T_{26}) = 26$$



# Example of Schedule Under TSO

- Is this schedule valid under TSO?
- How about this one, where initially
  - $R-TS(Q)=W-TS(Q)=0$
  - $TS(T_{27}) < TS(T_{28})$

$T_{27}$	$T_{28}$
read( $Q$ )	write( $Q$ )
write( $Q$ )	

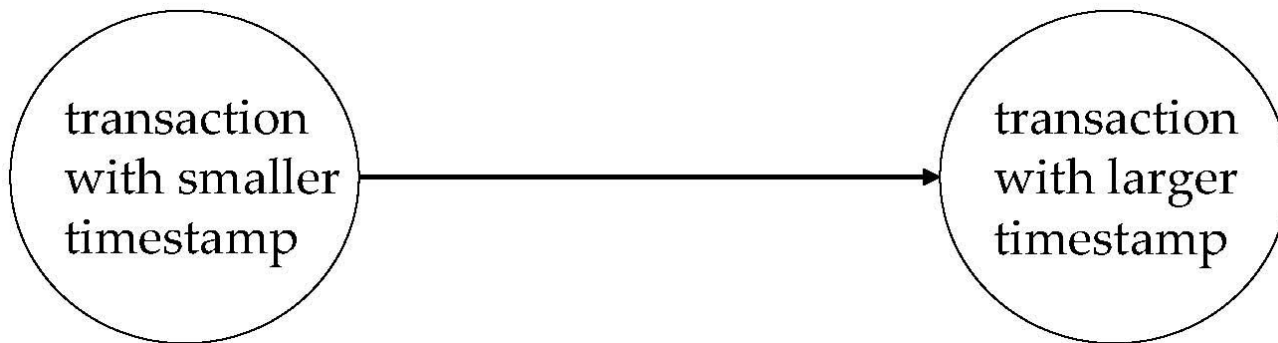
- The read( $Q$ ) operation of  $T_{27}$  succeeds, as does the write( $Q$ ) operation of  $T_{28}$ .
- When  $T_{27}$  attempts its write( $Q$ ) operation:
  - Because  $TS(T_{27}) < W\text{-timestamp}(Q)$ , since  $W\text{-timestamp}(Q) = TS(T_{28})$ . Thus, the write( $Q$ ) by  $T_{27}$  is rejected and transaction  $T_{27}$  must be rolled back.
- Any transaction  $T_j$  with  $TS(T_j) > TS(T_{28})$  must read the value of  $Q$  written by  $T_{28}$ , rather than the value that  $T_{27}$  is attempting to write





# Correctness of Timestamp-Ordering Protocol

- The timestamp-ordering protocol **guarantees serializability** since all the arcs in the precedence graph are of the form:



Thus, there will be no cycles in the precedence graph

- Timestamp protocol ensures freedom from deadlock as no transaction ever waits.
- But the schedule may not be cascade-free, and may not even be recoverable.



# Recoverability and Cascade Freedom

- Solution 1:
  - A transaction is structured such that its writes are all performed at the end of its processing
  - All writes of a transaction form an atomic action; no transaction may execute while a transaction is being written
  - A transaction that aborts is restarted with a new timestamp
- Solution 2:
  - Limited form of locking: wait for data to be committed before reading it
- Solution 3:
  - Use commit dependencies to ensure recoverability



# Thomas' Write Rule

- Modified version of the timestamp-ordering protocol in which obsolete **write** operations may be ignored under certain circumstances.
- When  $T_i$  attempts to write data item  $Q$ , if  $TS(T_i) < W\text{-timestamp}(Q)$ , then  $T_i$  is attempting to write an obsolete value of  $\{Q\}$ .
  - Rather than rolling back  $T_i$  as the timestamp ordering protocol would have done, this **{write}** operation can be **ignored**.
- Otherwise this protocol is the same as the timestamp ordering protocol.
- Thomas' Write Rule allows greater potential concurrency.
  - Allows some **view-serializable schedules** that are **not conflict-serializable**.



# Validation-Based Protocol

- Idea: can we use commit time as serialization order?
- To do so:
  - **Postpone writes to end** of transaction
  - Keep track of data items read/written by transaction
  - **Validation** performed at **commit time**, **detect** any **out-of-serialization order reads/writes**
- Also called as **optimistic concurrency control** since transaction executes fully in the hope that all will go well during validation



# Validation-Based Protocol

- Execution of transaction  $T_i$  is done in three phases.
  1. **Read and execution phase:** Transaction  $T_i$  **writes** only to **temporary local variables**
  2. **Validation phase:** Transaction  $T_i$  performs a "validation test" to determine if local variables can be written without **violating serializability**.
  3. **Write phase:** If  $T_i$  is validated, the updates are applied to the database; otherwise,  **$T_i$  is rolled back**.
- The three phases of concurrently executing transactions can be **interleaved**, but each transaction must go through the three phases in that order.
  - We assume **for simplicity** that the **validation and write phase** occur together, **atomically and serially**
    - I.e., only one transaction executes validation/write at a time.



## Validation-Based Protocol (Cont.)

- Each transaction  $T_i$  has 3 timestamps
  - **StartTS**( $T_i$ ) : the time when  $T_i$  started its execution
  - **ValidationTS**( $T_i$ ): the time when  $T_i$  entered its validation phase
  - **FinishTS**( $T_i$ ) : the time when  $T_i$  finished its write phase
- Validation tests use above timestamps and read/write sets to ensure that serializability order is determined by validation time
  - Thus,  $TS(T_i) = \text{ValidationTS}(T_i)$
- Validation-based protocol has been found to give greater degree of concurrency than locking/TSO if probability of conflicts is low.



# Validation Test for Transaction $T_j$

- **If** for all  $T_i$  with  $TS(T_i) < TS(T_j)$  either one of the following condition holds:
  - **finishTS( $T_i$ ) < startTS( $T_j$ )**
  - **startTS( $T_j$ ) < finishTS( $T_j$ ) < validationTS( $T_j$ )** and the set of data items written by  $T_i$  does not intersect with the set of data items read by  $T_j$ .
    - then **validation succeeds** and  $T_j$  can be committed.
- **Otherwise**, validation fails and  $T_j$  is aborted.
- Justification:
  - First condition applies when execution is not concurrent
    - The writes of  $T_j$  do not affect reads of  $T_i$  since they occur after  $T_i$  has finished its reads.
  - If the second condition holds, execution is concurrent,  $T_j$  does not **read any item written by  $T_i$** .



# Schedule Produced by Validation

- Example of schedule produced using validation Suppose that  $TS(T_{25}) < TS(T_{26})$ .

$T_{25}$	$T_{26}$
read( $B$ )	read( $B$ ) $B := B - 50$ read( $A$ ) $A := A + 50$
read( $A$ ) <validate> display( $A + B$ )	<validate> write( $B$ ) write( $A$ )

- The validation phase succeeds in the schedule.
- Note that the writes to the actual variables are performed only after the validation phase of  $T_{26}$ . Thus,  $T_{25}$  reads the old values of  $B$  and  $A$ , and this schedule is serializable.